NKind: a model checker for liveness property verification on Lustre programs

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- Model checking is getting greater concern in industrial field, like real-time reactive systems
- · Existing tools mostly concentrate on proving safety properties
- We present NKind, the first model checker for Lustre supporting the verification of liveness properties

• Lustre

- A synchronous dataflow language widely used in modeling reactive control systems
- Focus on data and represent them as infinite sequences of values, i.e. dataflows
- A Lustre program consist of nodes
 - · Intputs/outputs as interface as well as local dataflow
 - · Contain flow definitions
 - · Nodes work synchronously, i.e. at the same speed

```
node main() returns (counter2:int);
var
    counter1 : int;
let
    counter1 = 0 -> 1 + pre counter1;
    counter2 = 100 -> (pre counter2 - counter1);
tel;
```

counter1	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
counter2	100	99	97	94	90	85	79	72	64	55	45	34	22	9	-5

• L2SIA-WFR

- Extended algorithm of liveness-to-safety for handling infinite state space using abstraction
- Iterates and refines the abstraction by making use of spurious counterexamples
- Try to prove the absence of a lasso-shape counterexample as an invariant

Architecture

- Parser
 - Accepts an extension to core Lustre
- Simplifier
 - Automata, Array iterator, etc
- Controller
 - Dispatches the proving task
- Parallel engines
 - BMC
 - k-Induction
 - Invariant Generation
 - PDR

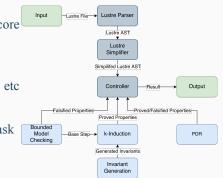


Figure 1: NKind architecture

Features

- Safety property
 - Attempts to prove that the given properties are invariants
 - Make use of the multiple engines mentioned before

```
node main() returns(x: int; y: int; f: bool);
var
e1: bool;
counter: int;
let
x = 20 -> (pre x) + 1;
y = 0 -> (pre y) + 2;
f = false -> pre x counter = 0 -> pre counter + 1;
e1 = (counter > 21) => f;
--$MATN;
--$PROPERTY Safe e1;
tel;
```

Figure 2: Specifying property with safety property



Figure 3: Safety verification result

Features

- Liveness property
 - Support properties which can have finite-time violations and will finally holds forever
 - Integrate L2SIA-WFR to PDR process

```
node main() returns(x: int; y: int; f: bool);
var
    e1: bool;
let
    x = 20 -> (pre x) + 1;
    y = 0 -> (pre y) + 2;
    f = false -> pre x     e1 = f;
    --%MAIN;
    --%PROPERTY Live e1;
tel;
```

Figure 4: Specifying property with liveness property

```
VALID PROPERTIES: [Liveness Property: e1] || Pdr || Step = 1 || Time = 2 s 482 ms
------
SUMMARY
------
VALID PROPERTIES: [Liveness Property: e1]
```

Figure 5: Liveness verification result

Features

- Liveness property
 - Structure referred to ic3ia
 - Monotonically strengthens the original transition system
 - Refine abstraction from spurious counterexamples

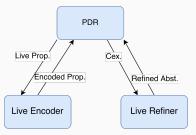


Figure 6: Embedding liveness extension to PDR process

Evaluation

• Safety benchmark from Kind (864)

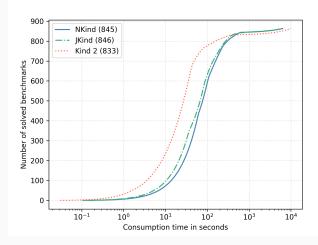


Figure 7: Verification results on safety property benchmark

Evaluation

• Liveness benchmark from ic3ia

Test name	ic3ia	NKind		
any-down-live	35.77	1.88		
parallel-live	51.68	12.61		
binary-live	0.09	0.86		
piecewise-live	0.11	1.13		
count-nested-live	0.37	1.87		
stabilize-live	5.16	3.99		
count-down-live	1.05	1.00		
swap-dec-live	1.53	2.73		
count-up-to-sym-live	5.84	1.11		
refine_disj_problem	Timeout	Timeout		

Table 1: Representative verification results on liveness property benchmark

Evaluation

• Liveness benchmark transformed from Kind (225)

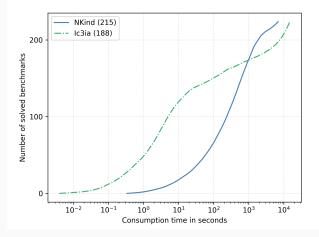


Figure 8: Verification results on more complex liveness property benchmark

- We present NKind, the first model checker for Lustre supporting the verification of liveness properties
- NKind has a rather competitive performance comparing to mainstream tools
- More in-depth optimizations for performance is left for future work

Thank you!